

Lecture 17

Boosting

Weak Learning

def. algorithm A weakly PAC learns concept class C if $\exists \gamma > 0$ st.

$\forall c \in C$ & \forall dists \mathcal{D} ,

given examples of c according to \mathcal{D}

A outputs h st. $\Pr_{\mathcal{D}} [h(x) \neq c(x)] \leq \frac{1}{2} - \frac{\gamma}{2}$

↑
advantage

Thm if C can be weakly PAC learned (on any \mathcal{D}) then

C can be (strongly) PAC learned.

Weak vs. Strong Learning

Def. Algorithm A weakly "PAC learns" concept class \mathcal{C}

if $\forall c \in \mathcal{C}$ & \forall dists \mathcal{D} $\exists \delta > 0$

$\forall \epsilon, \delta > 0$ ($\delta = \frac{1}{4}$ or $\frac{1}{n^2}$ doesn't affect)

with prob $\geq 1 - \delta$
given examples of c

A outputs h s.t. $\Pr_{\mathcal{D}} [h(x) \neq c(x)] \leq \frac{1}{2} - \frac{\delta}{2}$

\uparrow
advantage

It was conjectured that distribution free weak learning
was really weaker but surprise!

can "boost" a weak learner

Thm if \mathcal{C} can be weakly learned on
any dist \mathcal{D} then \mathcal{C} can be
(strongly) learned.

Applications

1) "Theoretical"

- Unif dist Algorithms for poly term DNF
weight w - poly threshold fctns

} low degree
alg doesn't
work well

∴ (Boosting + KM)

- Ave case vs. worst case

2) practical - Boosting
Freund-Schapire

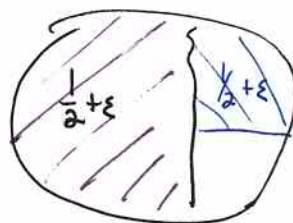
Good & Bad Ideas

- 1) simulate weak learner several times on
same distribution & take majority answer
-or-
best answer

gives better confidence

but doesn't reduce error, what if always get same answer?

- 2) filter out examples on which current hypothesis
does well & run weak learner on part where you
do badly.



Problem: given a new
example, how do you
know which section it
is in?

3) **Keep** some samples on which you are ok
always use **majority vote** on all previous hypotheses
to predict value of new samples

history: Schapire, Freund-Schapire, Impagliazzo-Servedio, Klivans

Filtering Procedures

- decide which samples to keep, which to throw out
- samples on which so far you guess correctly \leftarrow need for checking future hypotheses
- samples on which so far you guess incorrectly \leftarrow need to improve on these

The setting

- Given labeled examples
 $(x_1, f(x_1)), (x_2, f(x_2)), \dots$
 $x_i \in \mathcal{X}$
 $f \in \mathcal{C}$
- Given weak learning alg WL which weakly learns (advantage $\frac{\epsilon}{2}$) on any dist \mathcal{D}

Boosting Algorithm

Stage 0 (Initialize)

$$D_0 \leftarrow D$$

run WL on D_0 to generate (whp)

$$C_1 \text{ s.t. } \Pr_{D_0} [f(x) = C_1(x)] \geq \frac{1}{2} + \gamma/2$$

• For $i = 1 \dots T = O(\frac{1}{\gamma^2 \epsilon^2})$ stages, stage i : (can stop if Majority(C_1, \dots, C_i) correct on $\geq 1-\epsilon$ inputs)

(1) Construct D_i via "filtering procedure":

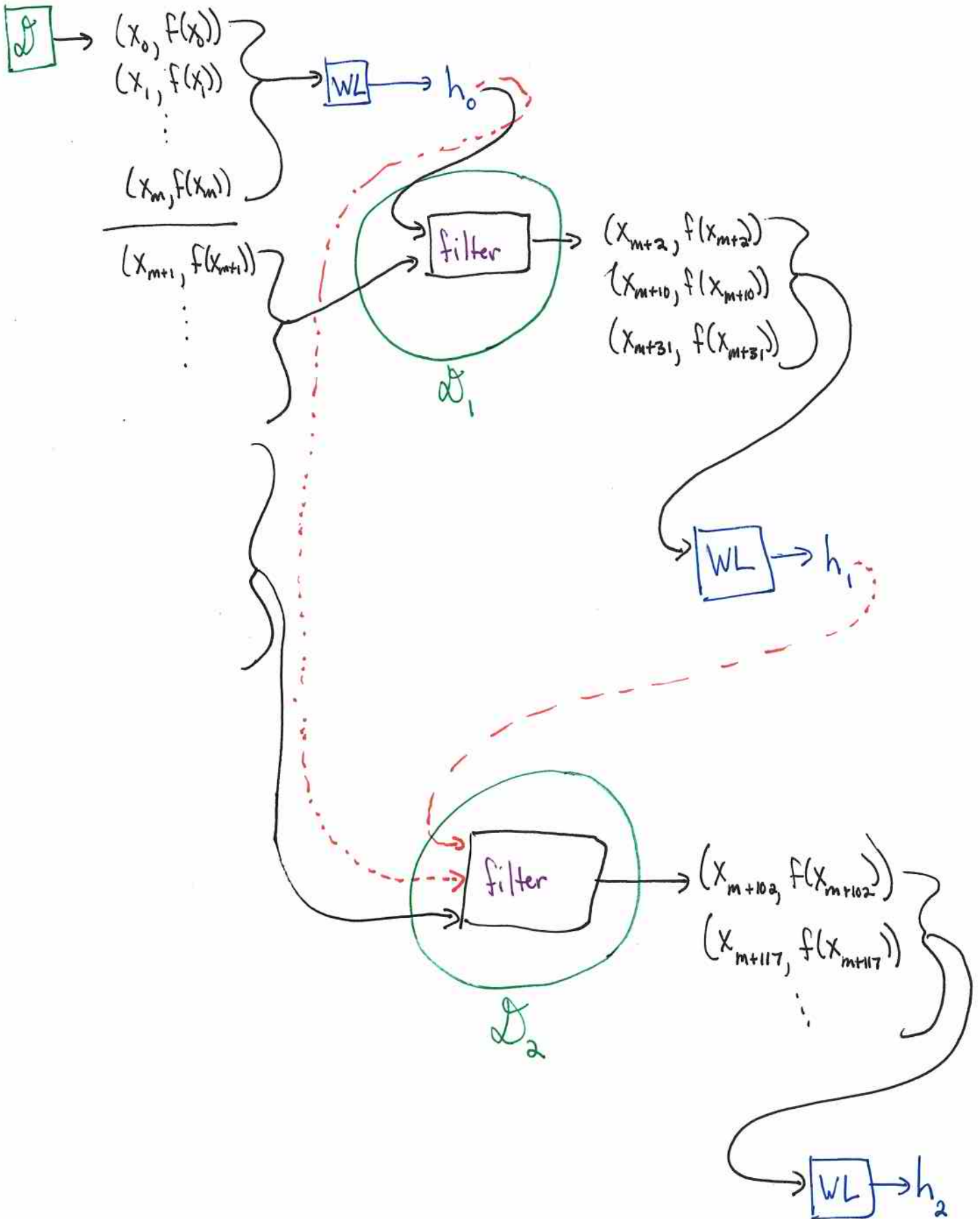
{ favor pts on which maj of $C_1 \dots C_i$ don't do well
but also keep some other points }

Will specify soon

(2) run WL on examples from D_i to output

$$C_{i+1} \text{ s.t. } \Pr_{D_i} [f(x) = C_{i+1}(x)] \geq \frac{1}{2} + \frac{\gamma}{2}$$

• output $C = \text{MAJ}(C_1 \dots C_T)$



Filtering procedure

Given new example $x, f(x)$ from example oracle

• if majority of $C_1 \dots C_i$ wrong, keep it
ie. $\geq \frac{i}{2}$

• if large majority right, then discard

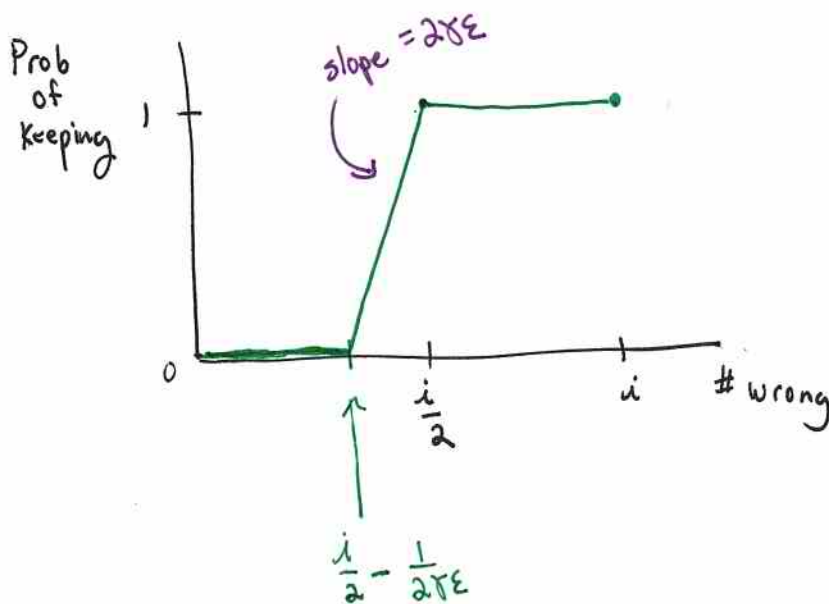
$$\text{ie. } \# \text{right} - \# \text{wrong} > \frac{1}{\gamma \epsilon}$$

$$\text{or } \# \text{wrong} \leq \frac{i}{2} - \frac{1}{2\gamma \epsilon}$$

• else $\# \text{right} - \# \text{wrong} = \frac{\alpha}{\gamma \epsilon}$ for $0 < \alpha < 1$

$$\# \text{wrong} - \# \text{right} = \frac{-\alpha}{\gamma \epsilon}$$

So keep with prob = $1 - \alpha$



Need to show:

1) Output is has nontrivial agreement with f

2) # samples needed not too bad

why could it be bad?
if throw out lots of samples, might
need to wait a long time before WL
can give an output, too many samples then
but if throw out too many samples then
you already have a good hypothesis!



will stop if $\text{Maj}(C_1 \dots C_i)$ correct on $\geq 1-\epsilon$ fraction
of inputs

o.w. $\text{Maj}(C_1 \dots C_i)$ incorrect on $> \epsilon$ fraction

so filtering procedure outputs
sample with prob $\geq \epsilon$

(+ in expectation, every $1/\epsilon$ samples
of \mathcal{D} at least one makes
it thru the filtering
system)

\Rightarrow filtering slows down sample
collection by $\leq O(1/\epsilon)$

So lets focus on ①

Notation

• $R_c(x) = \begin{cases} +1 & \text{if } f(x) = c(x) \\ -1 & \text{if } f(x) \neq c(x) \end{cases}$ "is c correct on x?"

• $N_i(x) = \sum_{1 \leq j \leq i} R_{c_j}(x)$ after iteration i , how many c's correct? (#right - #wrong)

• $M_i(x) = \begin{cases} 1 & \text{if } N_i(x) \leq 0 \\ 0 & \text{if } N_i(x) \geq \frac{1}{\epsilon} \\ 1 - \epsilon \cdot \delta \cdot N_i(x) & \text{o.w.} \end{cases}$ prob of keeping x in filtering (after stage i)
note - all "wrong" x included in M also some "right" x included

Note that new distribution on samples is proportional to M_i :

• $D_{M_i}(x) = \frac{M_i(x)}{\sum_x M_i(x)}$ distribution induced by M
note $D_{M_i}(x) = \mathcal{D}_i$

$\sum_x M_i(x)$ includes all "wrong" x but also x for which maj that suit are correct } upper bounds # wrong x

How correct are we wrt. D_{M_i} ?

• $Adv_c(M_i) = \sum_x R_c(x) M_i(x)$

• $\Pr_{x \in D_{M_i}} [c(x) = f(x)] = \frac{1}{2} + \frac{Adv_c(M_i)}{2 \cdot \sum_x M_i(x)}$
 $\frac{1}{2}$

"Advantage" of c on M
 $\sim \Pr[\text{correct}] - \Pr[\text{incorrect}]$
 $= 2 \cdot \Pr[\text{correct}] - 1$

Note:

$$\text{if } \sum M_i(x) \geq \epsilon 2^n$$

$$\text{Adv}_c(M_i) \cong \gamma \cdot \epsilon \cdot 2^n$$

convert claim about WL \Rightarrow claim about advantage
 i.e. if have \uparrow advantage on output of WL
 + still almost wrong on lots of inputs
 then new advantage is pretty good
 if not, then you are done

Begin Proof

For input x

$$\text{let } A_i(x) \leftarrow \sum_{0 \leq j \leq i-1} R_{c_{j+1}}(x) M_j(x)$$

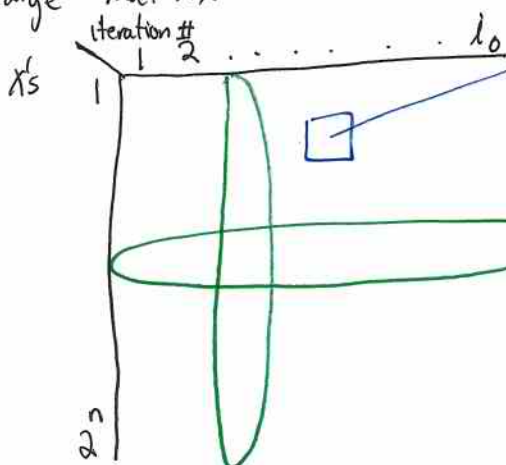
strange -
 indices don't match
 $c_1 \dots c_j$ define D_j
 but c_{j+1} learned from
 WL on D_j

Claim $A_i(x) \leq \frac{1}{\epsilon \gamma} + \frac{\epsilon \gamma}{2} \cdot i$

- bounds advantage per input
- only helps after $\frac{1}{\epsilon \gamma}$ rounds

Plan for use of claim:

Consider large matrix:



$(k, j+1)^{\text{th}}$ Entry: $R_{c_{j+1}}(x_k) M_j(x_k)$

$$x\text{'s row sum} = \sum_{0 \leq j \leq i_0} R_{c_{j+1}}(x) M_j(x) = A_i(x)$$

$$j^{\text{th}} \text{ col sum} = \sum_x R_{c_{j+1}}(x) M_j(x) = \text{Adv}_{c_{j+1}}(M_j) \leftarrow \geq \gamma \sum_x M_j(x)$$

else algorithm stops