# The Space "Just Above" BQP

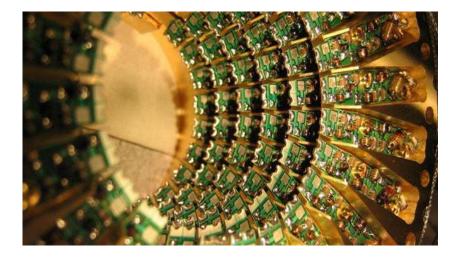
Adam Bouland

Based on joint work with Scott Aaronson, Joseph Fitzsimons and Mitchell Lee arXiv: 1412:6507 ITCS '16

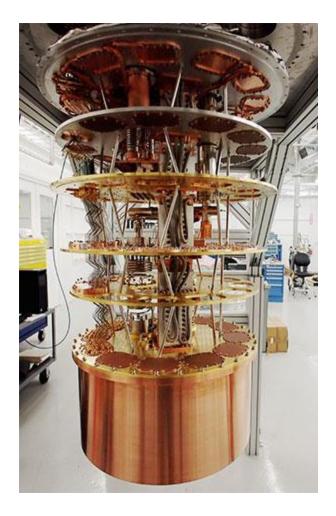




### **Quantum Computers**







### Quantum Computers...

#### **CAN** efficiently

• Factor integers [Shor]

### **CANNOT** efficiently

- Solve black-box NP-hard problems [BBBV]
  Searching N item list takes θ(N^1/2) time
- Solve black-box SZK-hard problems [Aaronson]

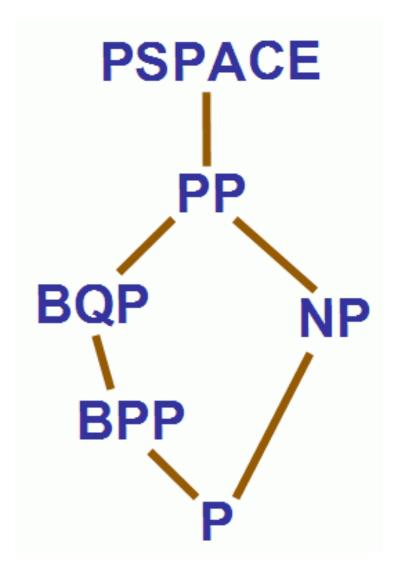


Image credit: Scott Aaronson

### **Quantum Mechanics**

- 1. State is vector  $v \in \mathbb{C}^d$   $||v||^2 = 1$
- 2. Unitary Evolution:  $v \to Uv$
- 3. Measurement

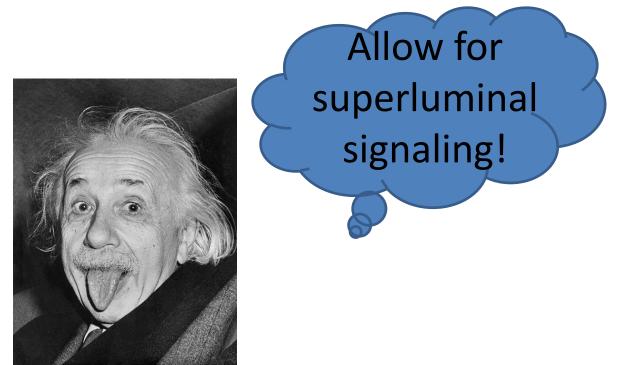
$$x \text{ w.p. } | < e_x, v > |^2$$
  
 $v \rightarrow e_x$ 

"Wavefunction Collapse"

### **Quantum Mechanics**

What happens to the power quantum computing if we perturb these axioms?

- Non-unitary evolution [Abrams-Lloyd], [Aaronson]
- Measurement based on p-norm for p!=2 [Aaronson]

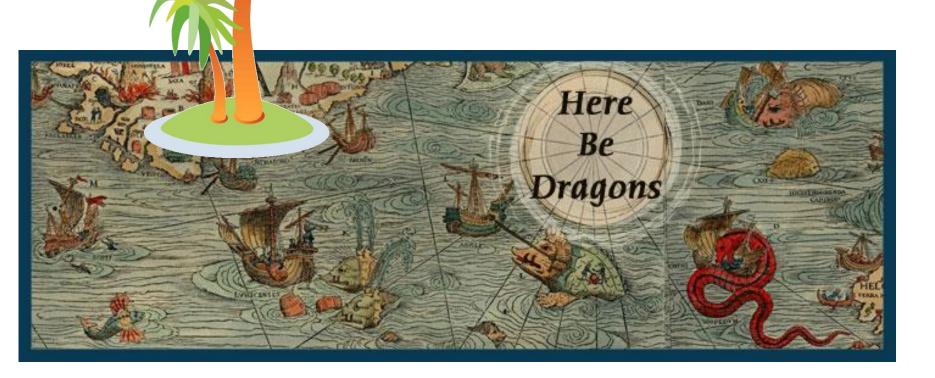


- Non-unitary evolution [Abrams-Lloyd], [Aaronson]
- Measurement based on p-norm for p!=2 [Aaronson]



- Non-unitary evolution [Abrams-L/ BOP -> PP Measurement based on p-norm for p: [Aaronson]





Challenge:

Is there *any* modification of QM that boosts the power of quantum computing to something SMALLER than PP?

> Yes (if you're careful)

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Is there *any* modification of QM that boosts the power of quantum computing to something SMALLER than <del>PP</del>-NP?

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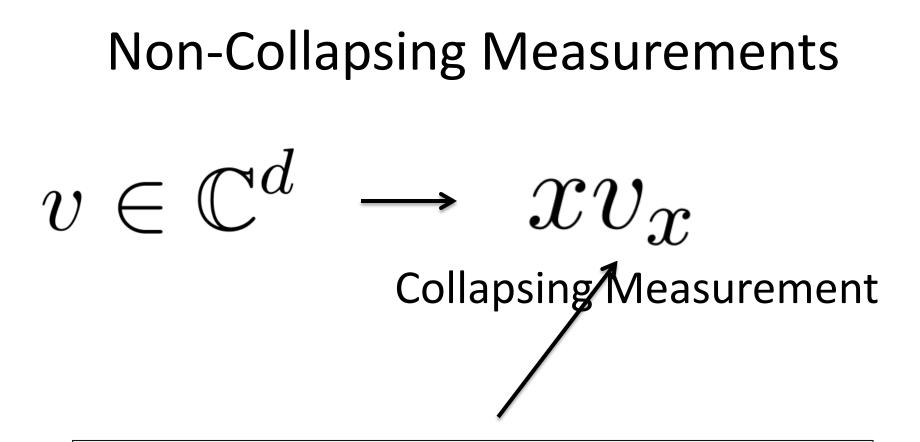
# **Non-Collapsing Measurements** $v \in \mathbb{C}^d$ Sample $x \text{ w.p. } | < e_x, v > |^2$ $v \to v$ "Wavefun Collapse"

#### **Non-Collapsing Measurements**

 $v \in \mathbb{C}^d$ x w.p.  $| < e_x, v > |^2$ 

x w.p.  $| < e_x, v > |^2$ 

v v



# Can measure same collapsed state multiple times

# Non-Collapsing Measurements

# CQP

"Collapse-free Quantum Polynomial time"

# naCQP

"non-adaptive CQP"

Quantum circuit must be non-adaptive to the non-collapsing measurement outcomes

**Non-Collapsing Measurements** 

# How powerful are these classes?

# A: naCQP is "just above" BQP

# Results

The class naCQP:

- Can solve SZK in poly-time
  - BQP cannot do this in black box manner
  - −  $\exists$  O such that naCQP^O  $\neq$  BQP^O
- Can search in O(N^1/3) time
- Search requires  $\Omega(N^1/4)$  time
  - −  $\exists$  O such that NP^O  $\not\subset$  naCQP^O
- In BPP^PP

### Summary

Property	BQP	naCQP
Contains SZK	Unknown	Yes
Contains $SZK^O \ \forall O$	No	Yes
Upper Bound for Search	$O(N^{1/2})$	$\tilde{O}(N^{1/3})$
Lower Bound for Search	$\Omega(N^{1/2})$	$\Omega(N^{1/4})$
Upper Bound	AWPP	BPP <sup>PP</sup>

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Upper Bound	AWPP	<b>B</b> PP <sup>PP</sup>	BPP <sup>PP</sup>

# Relation to Prior work

# Aaronson '05: QC with Hidden Variable Theories "DQP"

Imagine a hidden variable theory is true, and you "see" hidden variables of your system as it evolves

### Relation to Prior work

Property	BQP	naCQP	CQP	DQP
Contains SZK	Unknown	Yes	Yes	Yes
Contains $SZK^O \ \forall O$	No	Yes	Yes	Yes
Upper Bound for Search	$O(N^{1/2})$	$\tilde{O}(N^{1/3})$	$ ilde{O}(N^{1/3})$	$\tilde{O}(N^{1/3})$
Lower Bound for Search	$\Omega(N^{1/2})$	$\Omega(N^{1/4})$	$\Omega(1)$	Ω() 3)
Upper Bound	AWPP	BPPPP	BPP <sup>PP</sup>	ĔXP

# Don't bet on this model just yet!

- FTL Signaling (if adaptive)
- No notion of query complexity
- Can clone if circuit adaptive
  Perfect cloning-> #P [Bao B. Jordan '15]

– Imperfect cloning -> ???

### **Open Problems**

Property	BQP	naCQP	CQP
Contains SZK	Unknown	Yes	Yes
Contains $SZK^O \ \forall O$	No	Yes	Yes
Upper Bound for Search	$O(N^{1/2})$	$\tilde{O}(N^{1/3})$	$\tilde{O}(N^{1/3})$
Lower Bound for Search	$\Omega(N^{1/2})$	$\Omega(N^{1/4})$	$\Omega(1)$
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### Questions

2